

Lecture 8. Repeated Games

What happens if the same game is repeated? Can promises and threats about the future change the behavior?

1. Finitely repeated games

Definition: Given a static game G , the finitely repeated game $G(T)$ is a game in which G is played T times. For each $t \leq T$, the outcomes of the $t - 1$ preceding repetitions is known to the players before repetition t begins. Denote by π_i^t the payoff of player i from repetition t . The overall payoffs of player of the whole game is the sum of his payoffs of the T static games, i.e.

$$\pi_i^1 + \pi_i^2 + \dots + \pi_i^T = \sum_{t=1}^T \pi_i^t$$

To distinguish between the whole repeated game, G is also called "stage game".

Example: cooperation game (prisoner's dilemma), repeated twice.

	C_2	NC_2	
C_1	2 2	0 3	
NC_1	3 0	1 1	

This repeated game is a two-stage game with imperfect information (last lecture) \Rightarrow *SPE*

Second stage game: Only *NE* of the second stage is (NC_1, NC_2) , independently of what has happened in the first stage game.

First stage game: If in the second stage players play (NC_1, NC_2) independently of the result of first stage, (NC_1, NC_2) is the only equilibrium in the first stage, too.

\Rightarrow *SPE* of the twice repeated cooperation game is just twice the *NE* of the stage game.

Proposition: If the stage game G has a unique NE , then for any finite T the repeated game $G(T)$ has a unique subgame-perfect outcome: the NE of G is played in every stage.

This proposition does not hold if the stage game has more than one NE .

Example: stage game G , repeated twice

	L	M	R
U	1 1	5 0	0 0
C	0 5	4 4	0 0
B	0 0	0 0	3 3

Note: The stage game G has two NE in pure strategies: (U, L) and (B, R)

Take the following strategies of the players for the whole repeated game $G(2)$

row player:

1. Stage (Repetition): C
2. Stage (Repetition): $\begin{cases} B & \text{if 1.Stage resulted in } (C, M) \\ U & \text{if 1.Stage did not result in } (C, M) \end{cases}$

column player

1. Stage (Repetition) : M
2. Stage (Repetition) : $\begin{cases} R & \text{if 1.Stage resulted in } (C, M) \\ L & \text{if 1.Stage did not result in } (C, M) \end{cases}$

Claim: This strategy combination is a *SPE*.

Proof: i) We have to show that for any outcome of the first stage the second stage behavior is consistent with a *SPE*.

Take the first stage outcome (C, M) with payoffs $(4, 4)$. For this first stage outcome and the connected payoffs, the game and its payoffs become

	<i>L</i>		<i>M</i>		<i>R</i>	
<i>U</i>	1 + 4	1 + 4	5 + 4	0 + 4	0 + 4	0 + 4
<i>C</i>	0 + 4	5 + 4	4 + 4	4 + 4	0 + 4	0 + 4
<i>B</i>	0 + 4	0 + 4	0 + 4	0 + 4	3 + 4	3 + 4

For this static game, (B, R) is a *NE*.

Take any first stage outcome that is not (C, M) , and denote the connected payoffs by (x, y) . For this first stage outcome and the connected payoffs, the game and its payoffs become

	L		M		R	
U	$1 + x$	$1 + y$	$5 + x$	$0 + y$	$0 + x$	$0 + y$
C	$0 + x$	$5 + y$	$4 + x$	$4 + y$	$0 + x$	$0 + y$
B	$0 + x$	$0 + y$	$0 + x$	$0 + y$	$3 + x$	$3 + y$

For any (x, y) , (U, L) is a NE .

\Rightarrow second stage behavior is consistent with SPE - in fact any behavior that is an NE of the stage game G

ii) We have to show that for the proposed second stage actions, the first stage actions form a *NE* of the whole game.

Taken the second stage actions as described above into account, the payoffs are given by:

	<i>L</i>		<i>M</i>		<i>R</i>	
<i>U</i>	1 + 1	1 + 1	5 + 1	0 + 1	0 + 1	0 + 1
<i>C</i>	0 + 1	5 + 1	4 + 3	4 + 3	0 + 1	0 + 1
<i>B</i>	0 + 1	0 + 1	0 + 1	0 + 1	3 + 1	3 + 1

Note: For first stage outcomes (C, M) the second stage payoffs are $(3, 3)$, whereas for all other first stage outcomes the second stage payoffs are $(1, 1) \Rightarrow$

(C, M) in the first stage is part of an *SPE*. ■

The possibility to make the second stage behavior dependent on the outcome of the first stage allows for "non-NE" behavior in the first stage. But this possibility only exists when the second stage behavior is not uniquely determined, i.e. G has more than one *NE*.

2. Infinitely repeated games

Definition: Given a static game, G , $G(\infty, \delta)$ denotes the infinitely repeated game in which G is repeated forever. For each t the outcomes of the $t - 1$ preceding repetitions is known to the players before repetition t begins. Denote by π_i^t the payoff of player i from repetition t . The overall payoffs of player of the whole game is the sum of his payoffs of all static games discounted by δ , i.e.

$$\pi_i^1 + \delta \pi_i^2 + \delta^2 \pi_i^3 + \dots = \sum_{t=1}^{\infty} \delta^{t-1} \pi_i^t$$

Alternatively, one can also use the average discounted payoff per stage

$$(1 - \delta) \sum_{t=1}^{\infty} \delta^{t-1} \pi_i^t$$

Note that without discounting nothing prevents the payoff to go to ∞ .

Folk-theorem: Let (π_1, \dots, π_n) be the payoff from a *NE* of the stage game G , and let (x_1, \dots, x_n) be any other feasible payoff vector of the stage game G with $x_i > \pi_i$ for all $i = 1, \dots, n$. If δ is sufficiently close to one, then there exists a *SPE* of the infinitely repeated game $G(\infty, \delta)$ that achieves (x_1, \dots, x_n) as average payoff.

"In an infinitely repeated game anything that is individually rational can be supported by a *SPE*."

Example: Cooperation game

	C_2	NC_2	
C_1	2 2	0 3	
NC_1	3 0	1 1	